## STACISS APPLICATIONS

## INFIX NOTHTION

$\star$ Infix notation is the common arithmetic and logical formula notation, in which operators are written infix-style between the operands they act on
$\star$ E.g. A + B

## POSTFIX NOTATION

$\star$ In Postfix notation, the operator comes after the Operand.
$\star$ For example, the Infix expression $\bar{A}+\mathbf{B}$ will be written as $\mathbf{A B}+$ in its Postfix Notation.
$\star$ Postfix is also called 'Reverse Polish Notation'

## PREFIX NOTATION

$\star$ In Prefix notation, the operator comes before the operand.
$\star$ The Infix expression $\boldsymbol{A}+\mathbf{B}$ will be written as $\mathbf{+} \boldsymbol{A} \mathbf{B}$ in its Prefix

Notation.

* Prefix is also called 'Polish Notation'


## BUILDING AN ARITHMETIC EXPRESSION

## Postfix Expression String Processing

> Assume 1-digit integer operands, the binary operators $+-\star /$ only, and the string to be evaluated is properly formed

Rules for processing the postfix string:
Starting from the left hand end, inspect each character of the string

1. if it's an operand - push it on the stack
2. if it's an operator - remove the top 2 operands from the stack, perform the indicated operation, and push the result on the stack

An Example: $3 *(4+5) / 2 \rightarrow 345+* 2 / \rightarrow 13$
Remaining Postfix String int Stack (top $\rightarrow$ ) Rule Used

| 345+*2/ | empty |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| $45+* 2 /$ | 3 |  |  | 1 |
| $5+* 2 /$ | 3 | 4 |  | 1 |
| +*2/ | 3 | 4 | 5 | 1 |
| * $2 /$ | 3 | 9 |  | 2 |
| $2 /$ | 27 |  |  | 2 |
| / | 27 | 2 |  | 1 |
| null | 13 |  |  | 2 |

## CONVERSION FROM INFIX TO POSTFIX HLGORITHM

Stepl
*Scan the Infix expression from left to right for tokens
(Operators, Operands \& Parentheses) and perform the steps
2 to 5 for each token in the Expression

## HLCORITHM

Step2
$\star$ If token is operand, Append it in postfix expression

Step3
$\star$ If token is a left parentheses "(", push it in stack.

## ALGORITHM

Step4
$\star$ If token is an operator,
>Pop all the operators which are of higher or equal precedence then the incoming token and append them (in the same order) to the output Expression.
>After popping out all such operators, push the new token on stack.

## ALGORITHM

Step5
*If ")" right parentheses is found,
>Pop all the operators from the Stack and append them to Output String, till you encounter the Opening Parenthesis "(".
>Pop the left parenthesis but don't append it to the output string (Postfix notation does not have brackets).

## HLCORITHM

Step6
*When all tokens of Infix expression have been scanned. Pop
all the elements from the stack and append them to the
Output String.
$\star$ The Output string is the Corresponding Postfix Notation.

## EXAMPLE

| An Example: $7-(2 * 3+5) *(8-4 / 2) \rightarrow 723 * 5+842 /-*-$ |  |  |  |
| :---: | :---: | :---: | :---: |
| Remaining Infix String | char Stack | Postfix String | Rule Used |
| $7-(2 * 3+5) *(8-4 / 2)$ | empty | null |  |
| $-(2 * 3+5) *(8-4 / 2)$ | empty | 7 | 1 |
| $(2 * 3+5) *(8-4 / 2)$ | - | 7 | 3 |
| $2 * 3+5) *(8-4 / 2)$ | - 1 | 7 | 2 |
| * $3+5)$ * (8-4/2) | - 1 | 72 | 1 |
| $3+5) *(8-4 / 2)$ | - ${ }^{*}$ | 72 | 3 |
| $+5)$ * (8-4/2) | - (* | 723 | 3 |
| 5) * $(8-4 / 2)$ | - ${ }^{+}$ | 723* | 3 |
| $) *(8-4 / 2)$ | - ${ }^{+}$ | $723 * 5$ | 1 |
| * (8-4/2) | - | 723*5+ | 4 |
| (8-4/2) | -* | 723*5+ | 3 |
| 8-4/2) | -* ( | 723*5+ | 2 |
| -4/2) | -* ( | $723 * 5+8$ | 1 |
| 4/2) | -* $(-$ | $723 * 5+8$ | 3 |
| /2) | -* $(-$ | $723 * 5+84$ | 1 |
| 2) | -* $(-1$ | $723 * 5+84$ | 3 |
| ) | -* (-) | $723 * 5+842$ | 1 |
| (Aull@@ | empty | $723 * 5+842 /-*-$ | $4 \& 5$ |

Example A* (B + C * D) + E becomes A B C D* +*E +
current symbol operator stack postfix string

| A |  | A |
| :---: | :---: | :---: |
| * | * | A |
| C | * | A |
| B | * | A B |
| $+$ | ${ }^{*}$ C + | A B |
| C | * + | ABC |
| * | * ${ }^{(+*}$ | $A B C$ |
| D | * ${ }^{(+*}$ | $A B C D$ |
| ) | * | $\mathrm{ABCD}^{++}$ |
| $+$ | $+$ | $A B C D+*$ |
| E | $+$ | $A B C D+{ }^{+}+\mathrm{E}$ |
|  |  | $A B C D *+*$ |

## EXAMPLE

$\star$ Let the incoming the Infix expression be:

$$
\mathbf{A} *(\mathbf{B}+\mathbf{C})-\mathbf{D} / \mathbf{E}
$$

Stage 1: Stack is empty and we only have the Infix
Expression.


## EXAMPLE

## Stage 2

$\star$ The first token is Operand $\bar{A}$ Operands are Appended to the
Output as it is.


InFix Notation:

* (B+C) - D / E

PostFix Notation:
A

## EXAMPLE

## Stage 3

$\star$ Next token is * Since Stack is empty (top==NULL) it is pushed into the Stack


InFix Notation:
(B+C)-D / E

PostFix Notation:
A

## EXAMPLE

## Stage 4

* Next token is (the precedence of open-parenthesis, when it is to go inside, is maximum.
$\star$ But when another operator is to come on the top of "( ${ }^{6}$ then its
precedence is least.


InFix Notation:
B+C) - D / E

PostFix Notation:
A

## EXAMPLE

## Stage 5

$\star$ Next token, B is an operand which will go to the Output expression
as it is


## EXAMPLE

## Stage 6

$\star$ Next token, + is operator, We consider the precedence of top element
in the Stack, '(‘. The outgoing precedence of open parenthesis is the
least (refer point 4. Above). So + gets pushed into the Stack


## EXAMPLE

Stage 7

* Next token, C, is appended to the output



## EXAMPLE

## Stage 8

$\star$ Next token ), means that pop all the elements from Stack and append
them to the output expression till we read an opening parenthesis.


## EXAMPLE

## Stage 9

$\star$ Next token, -, is an operator. The precedence of operator on the top of Stack '*' is more than that of Minus. So we pop multiply and append it to output expression. Then push minus in the Stack.


InFix Notation:
D / E

PostFix Notation:
ABC + *

## EXAMPLE

Stage 10
$\star$ Next, Operand 'D' gets appended to the output.


InFix Notation:
/E

PostFix Notation:
ABC + * D

## EXAMPLE

## Stage 11

* Next, we will insert the division operator into the Stack because its precedence is more than that of minus.



## EXAMPLE

## Stage 12

* The last token, E , is an operand, so we insert it to the output

Expression as it is.


## EXAMPLE

## Stage 13

*The input Expression is complete now. So we pop the Stack and Append it to the Output Expression as we pop it.


## EXAMPLE 2

$(((A+B) *(C-E)) /(F+G))$

- stack: <empty>
- output: []


## EXHMPLE 2

$((A+B) *(C-E)) /(F+G))$

- stack: (
" output: []


## EXAMPLE 2

$$
(A+B) *(C-E)) /(F+G))
$$

- stack: ( (
" output: []


## EXAMPLE 2

$A+B) *(C-E)) /(F+G))$

- stack: ( ( (
" output: []


## EXAMPLE 2

$+B) *(C-E)) /(F+G))$

- stack: ( (
" output: [A]


## EXAMPLE 2

B)*(C-E))/(F+G))

- stack: ( ( ( +
- output: [A]


## EXHMPLE 2

$$
) *(C-E)) /(F+G))
$$

- stack: ( ( ( +
- output: [A B]


## EXHMPLE 2

* (C-E))/(F+G))
- stack: ( (
- output: [A B + ]


## EXAMPLE 2

$$
(C-E)) /(F+G))
$$

- stack: ( (*
- output: [A B + ]


## EXHMPLE 2

$C-E) /(F+G))$

- stack: ( ( ${ }^{( }$
- output: [A B + ]


## EXHMPLE 2

$-E) /(F+G))$

- stack: ( ( ${ }^{( }$
- output: [A B + C ]


## EXHMPLE 2

E) $) /(F+G))$

- stack: ( ( *
- output: [A B + C ]


## EXHMPLE 2

)) $/(F+G))$

- stack: ( ( * (-
- output: [A B + C E ]


## EXHMPLE 2

$) /(F+G))$

- stack: ( (*
- output: [A B + C E - ]


## EXHMPLE 2

$/(F+G))$

- stack: (
- output: [A B + C E - * ]


## EXHMPLE 2

$(F+G))$

- stack: (/
- output: [A B + C E - * ]


## EXAMPLE 2

F + G ) )

- stack: (/ (
- output: [A B + C E - * ]


## EXHMPLE 2

+ G ) )
- stack: ( / (
- output: [A B + C E - * F ]


## EXAMPLE 2

G ) )

- stack: ( / (+
- output: [A B + C E - * F ]


## EXAMPLE 2

))

- stack: ( / (+
- output: [A B + C E - * F G ]


## EXHMPLE 2

)

- stack: (/
- output: [A B + C E - * F G + ]


## EXAMPLE 2

- stack: <empty>
- output: [A B + C E - * F G + / ]

